

CS 42—Inheritance and subtyping

Tuesday, December 4, 2018

Recall: object-oriented terminology

interface	<i>what</i> an object can do
type	a description of an object's interface
subtype	a type that extends the interface of another type (its supertype)
implementation	<i>how</i> an object does its thing
class	a description of an object's implementation
subclass	a class that extends the implementation of another class (its superclass)

Recall: Inheritance

Inheritance enables code reuse in two ways:

- Reuse for the benefit of providers: The provider of a new class can define that class by explaining how it is different from an existing class. (This kind of reuse is subclassing.)
- Reuse for the benefit of clients: The client of a type can write code that can be used with multiple implementations of that type. (This kind of reuse is subtyping.)

In most OO languages, when you use inheritance, you define a subclass *and* a subtype.

Good programming practice: use inheritance only if the existing class and the new class have an *is-a* relationship. Otherwise, it's probably better for the new class to contain a field whose type is the existing class.

Types in Java

The **declared type** of a variable is the type that appears in the code, to the left of the variable, when it is declared. For example:

```
Dog buddy; // buddy's declared type is Dog
Animal buddy; // buddy's declared type is Animal
```

The value of variable can always be used anywhere that type is expected, e.g., as an argument to a function, as a return value, etc.

Subtyping

Subtyping creates an “is-a” relationship: an instance of the subtype “is” an instance of the supertype. Here is how we can create is-a relationships (i.e., subtypes) in Java:

- Implementing an interface establishes an is-a relationship.
- Extending an interface establishes an is-a relationship.
- Extending a class establishes an is-a relationship.

Also, every variable “is an” instance of its declared type.

The “is-a” relationship is **transitive**: if A “is a” B and B “is a” C, then A “is a” C.

Actual types, subtyping, and substitutability

The **actual type** of a variable must match or be a subtype of the declared type. In other words, the actual object must support a superset of the methods described by its type. That way, we can be sure that every method call on the object is valid. As a result, subtyping allows us to substitute an instance of a subtype for an instance of a supertype because we know that the subtype supports all the operations in the supertype. In other words, if A is a subtype of B, then an instance of A can replace an instance of B in any situation that calls for a B. This idea is often referred to as the “Liskov substitution principle”, named after the researcher, Barbara Liskov, who initially introduced it.

Declared type *vs* actual type, in Java

When we compile a program, the type checker looks at the declared type (not the value) of an object to see *whether* the program’s method calls are legal.

When we run a program, Java uses the actual object (not the declared type of the object) to choose *which* method to run.

For example, consider the following code:

```
Dog spot = new FrenchPoodle("Spot", 4, 99);
spot.sayHello();
```

(where `FrenchPoodle` inherits from `Dog`). `Dog` is the declared type, and `FrenchPoodle` is the actual type. When we compile the program, the type checker will make sure that the `Dog` type defines a `sayHello` method. When we run the program, Java will call the version of `sayHello` that is defined in the `FrenchPoodle` class.

CS 42—Graphs

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Graph definitions and terminology

A **graph** contains **nodes** (also called **vertices**) and **edges**. An edge connects two nodes.

We can use graphs to represent relationships: there is a relationship between node A and node B if there is an edge between A and B.

In an **undirected graph**, relationships are mutual.

In a **directed graph**, relationships are one-way, from the **source** to the **destination**.

In a **weighted graph**, relationships have associated information (e.g., cost or “weight”).

An edge is **adjacent** to A if that edge emanates from A.

Node B is **adjacent** to node A if there is an edge from A to B.

The **neighbors** of A are all the nodes adjacent to A.

A **path** is a sequence of edges between adjacent nodes.

Node B is **reachable** from node A if there is a path from A to B.

In a **complete graph**, there exists an edge between each pair of nodes.

In a **connected graph**, there exists a path between each pair of nodes.

A **sparse** graph has few edges, relative to the maximum amount it can have.

A **dense** graph has many edges, relative to the maximum amount it can have.

In an **acyclic** graph, there are a finite number of paths between any two nodes.

In a **cyclic** graph, there may be an infinite number of paths between two nodes.

Designing and implementing a new data structure

The **interface** describes *what* a data structure can do (e.g., its operations). The interface is a promise from the provider of the data structure to the user of the data structure.

The **implementation** describes *how* the data structure works (e.g., how the data are stored / organized and which algorithms are used to provide the operations). The implementation makes good on the promise of the interface.

It should be possible to replace the implementation without modifying the interface.