

Final exam

- In class, Tuesday 12/18 9am–noon

In this room

If you have another final at that time, then Monday 12/19 2–5pm

Exam target: 75 minutes

You can bring 2, double-sided pages of notes

- Cumulative, but will focus a bit more on recent topics

OOP • dynamic programming • summations • graphs

nope: ~~Turing machines • proofs about languages • sorting~~

Poster session!

1	Circuits
2	Finite-state machines
3	How a Hmmm program runs <i>registers, RAM, stack, etc.</i>
4	Linear data structures <i>lists, arrays, stacks, queues</i>
5	The building blocks of functional programming <i>map, filter, foldl, foldr, list comprehensions</i>
6	Recursion and use-it-or-lose-it
7	Analysis techniques <i>summations, recurrence relations, asymptotic complexity</i>
8	Python namespaces
9	Dynamic programming <i>with tabulation</i>
10	Object-oriented programming <i>Python vs Java</i>
11	Trees and graphs <i>Properties and algorithms</i>

Brainstorm

Idea generators: call out things big and small related to this topic

Scribes: write down the ideas

You can switch roles!

Goal: quantity, not quality

we'll deal with quality next...





- Nothing is too big or too small
- **No judgements:** Nothing is better than anything else

Organize

- How might you start to arrange / organize / explain this topic?
Can you make sub-categories?
- What's most important?
- What's "good to know", but not that important?
- What's unimportant or unrelated?
- What do you feel you can explain clearly?
- What is not yet clear?

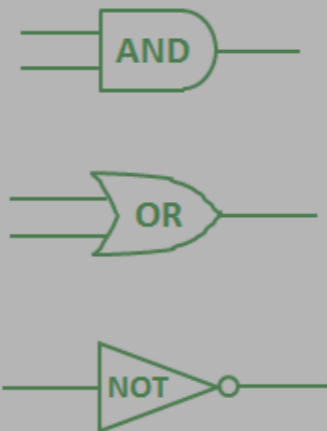
Posters!

13 myths about
functional programming
and one important truth

- 1 
- 2 
- 3 
- 4 

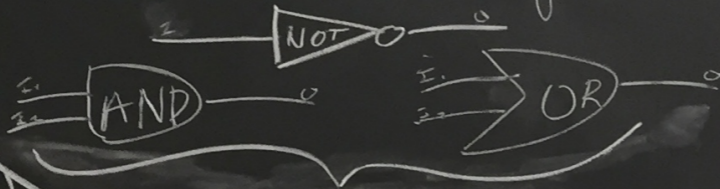
...

All My Circuits
*an exciting new
Netflix show*



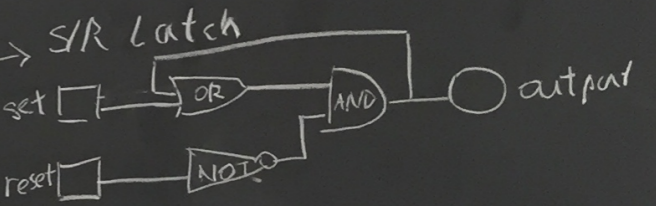
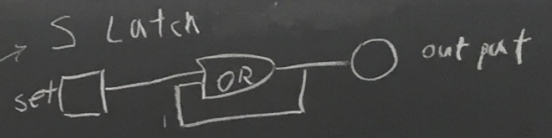
How to care for and
feed a Python

① Latch on to the Logic

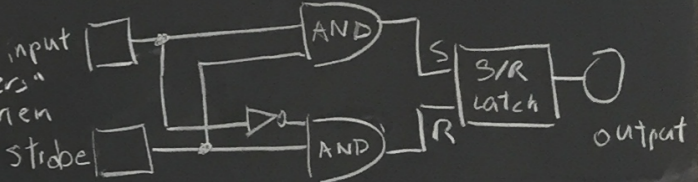


PUT IT TOGETHER TO MAKE PARTS OF COMPUTERS

once set = 1, remembers this input
S Latch with reset button



Pass through → D Latch
When strobe input is on, "remembers" last value when strobe off



Good Circuits

- Correct
- Simple, easy to read
- Efficient

FUN!

EVERYONE THINKS YOU NEED ALL 3 TO MAKE A COMPUTER, BUT WE ONLY USED 2!

COMPUTER SCIENTISTS HATE H.M!

IN	OUT
1	0
0	1

NOT

IN ₁	IN ₂	OUT
0	0	0
0	1	1
1	0	1
1	1	1

OR

IN ₁	IN ₂	OUT
0	0	0
0	1	0
1	0	0
1	1	1

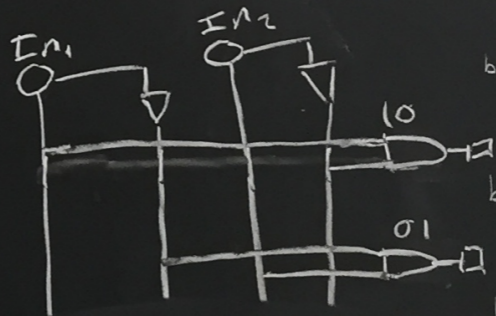
AND

Truth Table

in ₁	in ₂	out
0	0	0
1	0	1
0	1	1
1	1	0

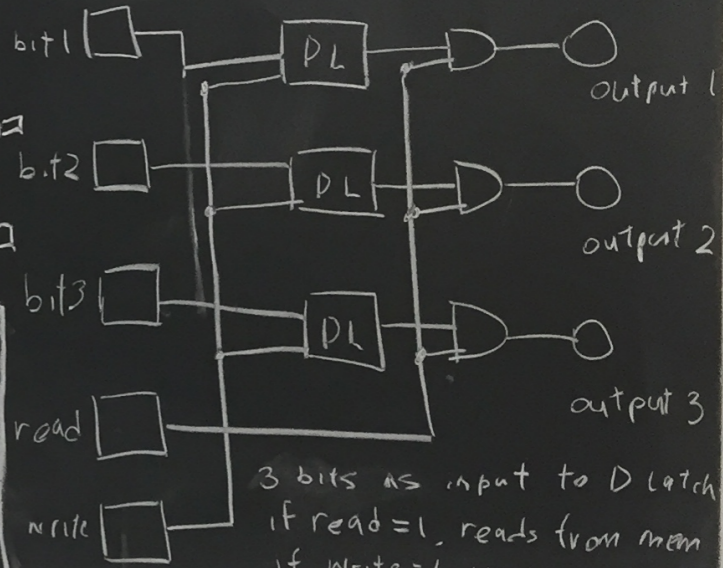
XOR

Minterm Expansion



Each output "i" gets an and gate connected to each input directly or through a not gate, and then goes to the output.

3-Bit Memory

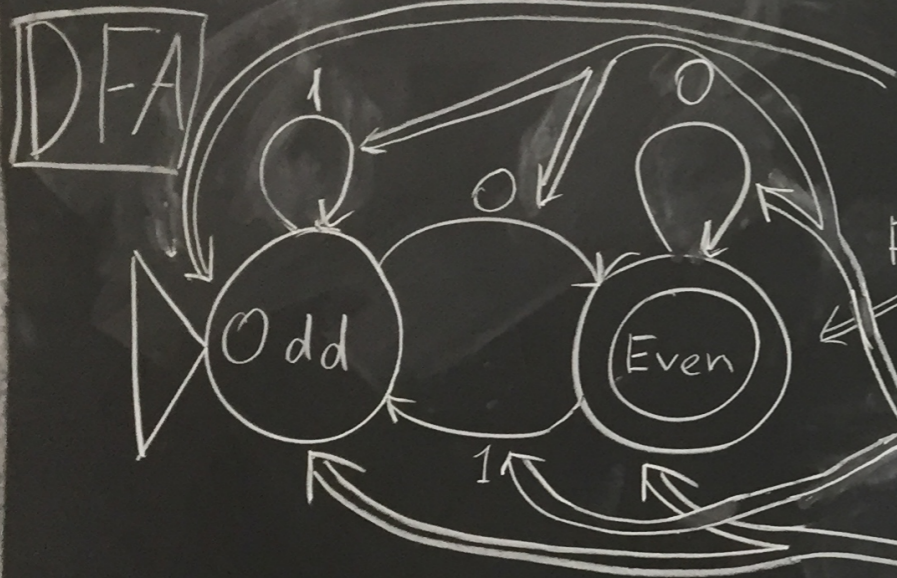


3 bits as input to D Latch
if read = 1, reads from mem
if write = 1, writes input to mem

ON
MS

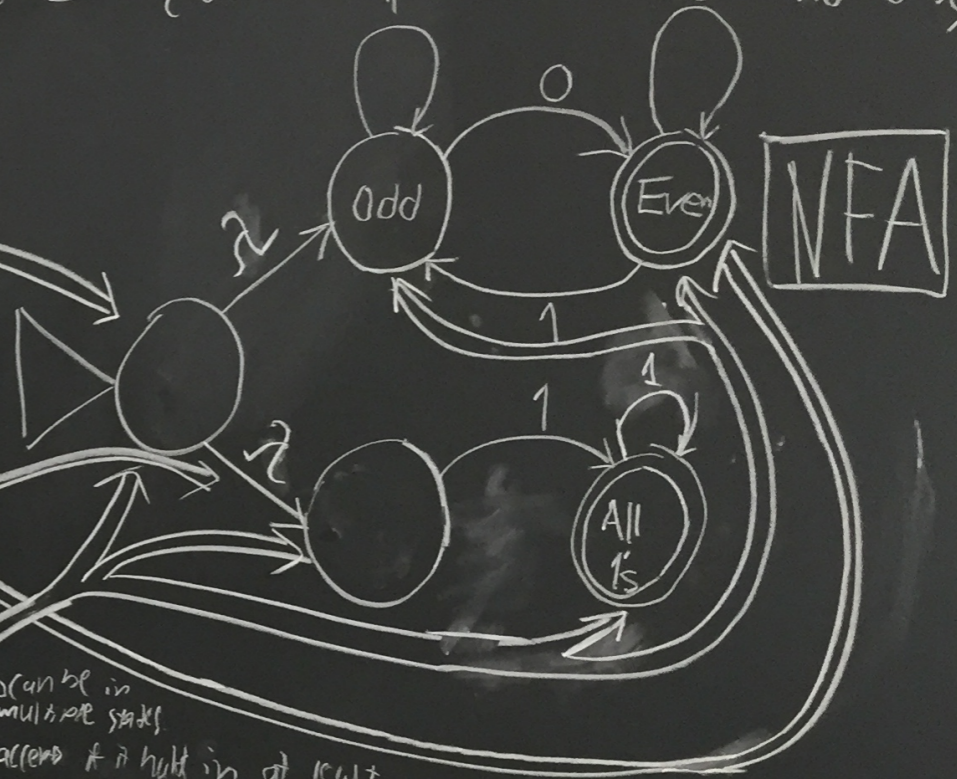
2. Finite State Machine

Language Input $L = [w \mid w \text{ is even}]$ \Rightarrow $L = \{w \mid w \text{ is even or } w \text{ contains no } 0\text{'s}\}$



DFA

Initial state
Final/Accepting state
Transition function
State



NFA

- Each state has exactly one transition for each character in the alphabet
- at any time it is in one state at a time
- Accepts if it halts in an accepting state

can be in multiple states
accepts if it halts in at least one accepting state

$NFAs \equiv DFAs \equiv REs$
Kleene's Theorem

Turing Machines \leftarrow All Decision Problems



OMG - 4 tips for HMMM that'll clear your skin

CPU

Fetch-Execute Cycle:

1. Retrieve the instruction at the program counter
 2. Execute the instruction
 3. Increment the program counter
- This is how HMMM executes a program

Instruction types

- Computations (add, subtract, etc registers)
- GOTOs (jump & various jumps)
- Memory management (move values between registers & RAM)

- 00 Set r15 Q # set stack pointer
- 01 Read r1 # read a val
- 02 Call r14 # go to func
- 03 write r15
- 0N Push r1
- 0N+1 Push r14 # save
- 0N+2 Do stuff
- 0N+? Jump 0N something
- 0N+? pop r14
- 0X pop r1
- 0Y copy r13 r1
- 0Z jump r14

Memory

- r₁₃: return value of function
- r₁₄: return line (where the program counter goes after function ends)
- r₁₅: stack pointer (points to top of stack)

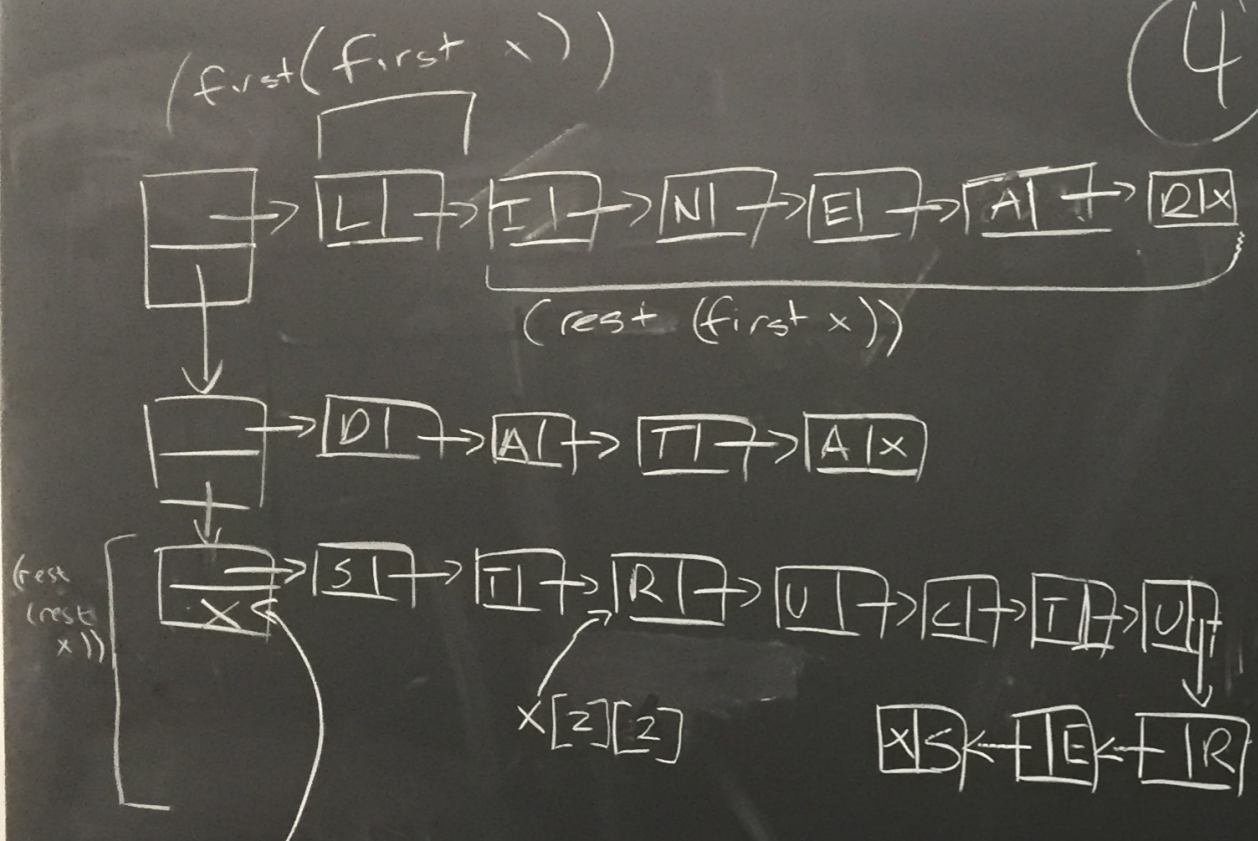
STACK: LIFO; stores variables throughout program outside of registers.

- Push: Add a value to top of stack
- Pop: Remove most recent value from stack
- Stack builds to higher RAM addresses

Other important skills

- Tracing code (given code, say what it does)

4



$(\text{first } x) = x[0]$

$(\text{rest } x) = x[1:]$

Python

lists: $[-, -, \dots, -]$ (mutable)
 tuples: $(-, -, \dots, -)$ (immutable)

$[x * 2 \text{ for } x \text{ in range}(10) \text{ if } x > 5]$
 map element list filter

Java

ArrayList<Object>
 int[]

Arrays: pros: search/access
 cons: resizing
 Better for queues (you have to grab the last element)

Empty List (Null)

Linked List (linked)

pros: insertion, deletion, first element
 cons: search/access
 haha

great for stacks

Stacks

pop - remove top element
 push - add element to top

Do string

Functional Programming!

PROS!

Referential Transparency

- same input \rightarrow same output
(replace variables with values)

- No side effects
(changes in state ie printing to screen)

\rightarrow Easier debugging

Order of execution

is irrelevant

Not as much memory

CONS

CONS

No side-effects

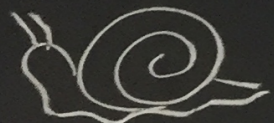
more abstract

harder to understand (at first)

redundant

of a list)

1. Base Case
2. "It's one piece"
3. What is the solution with it?
4. What is the solution without it?
5. How do we combine them?



redundant

Functional Racket

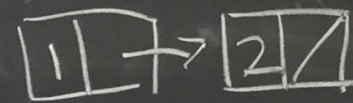
Imperative

Interface

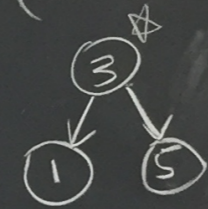
Implementation

Data Structures - Function
List

(define f(x) (+ x 1))



(cons 1 '(2))



tree

Anonymous functions

(define adder (lambda (n) (+ n x)))

HOFs - map
filter
fold

(map adder 1) L

(filter (lambda (n) (= 0 (% n 2))))

(fold + 0 L)
(1+2)+3

"Variables"

(constant values)

(let ([CS 42])
body expr)

Performance of a program
have finite resources
depends on situation

time, memory
calls, steps

RECURRENCE RELATIONS

$T(N)$ = cost given input size N

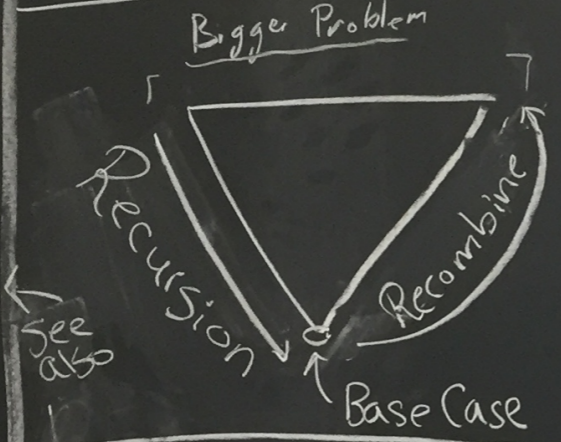
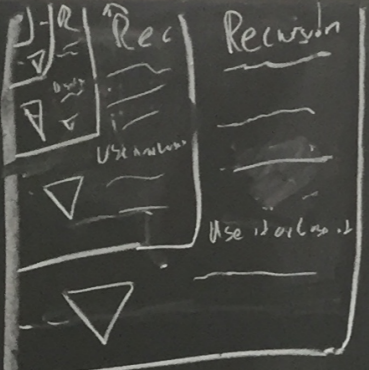
define (sum n)
(if $= n 0$)
0
(+ n
(sum
(- $n 1$)))

cost metric =
recursive calls
 $T(0) = 0$
 $T(1) = 1 + T(0)$
 $T(2) = 1 + T(1)$
 $T(N) = N$
which is $O(N)$

Used to find
to find time complexity
of recursive functions

tractibly
(N^N) $O(2^N)$
($N!$)

$T(N)$ can
be recursive



REPRESSED MEMORIES

- (1) Align
- (2) Make-change
- (3) Car pne
- (4) Tree traversals
- (5) Word difference (Edit Distance.py)

Recursion

- Solving a problem by breaking it into smaller problems (ex: sum of a list)

Use it or Lose it

- A specific type of recursion
- used by making a template:

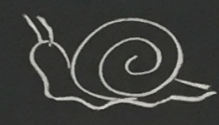
1. Base Case
2. It's one piece
3. What is the solution with it?
4. What is the solution without it?
5. How do we combine them?

Benefits

- Intuitive
- Gets Right answer
- Useful for inductive structures

Drawbacks

- Slow
- Inefficient, redundant Computations



Do string Fun =

PROS!

Referential Transparency

- same input \rightarrow same output
- (replace variables w/ vals)
- No side effects (changes in state ie printing to screen)
- \rightarrow Easier debugging

- Order of execution is irrelevant

- Not as much memory

CONS

- No side-effects
- more abstract
- harder to understand (at first)

Analysis Techniques

What: Measuring the performance of a program
 why: Real world programs have finite resources

CONCEPTS Cost metric: some metric of cost, depends on situation
 Empirical: measured Ex. time, memory
 theoretical: ↓ Ex. calls, steps

SUMMATION

$T(N)$ = cost given input size N

for j in range (N)

$$\sum_{j=0}^{N-1} k = N \cdot k \leftarrow \text{if } k \text{ is constant}$$

nested loop: $\sum_{\text{outer loop}} \sum_{\text{inner loop}} k \in O(N^2)$

*USE for loops

no problem
 $O(\log_2 N)$, $O(1)$
 $O(\sqrt{N})$

tractable
 $O(N)$
 $O(N \log N)$

Intractable
 $O(N^N)$, $O(2^N)$
 $O(N!)$

$\hookrightarrow T(N)$ can be recursive

RECURRENCE RELATIONS

$T(N)$ = cost given input size N

(define (sum n)

(if (= n 0)

0

(+ n

(sum

(- n 1))))

cost metric = recursive calls

$T(0) = 0$

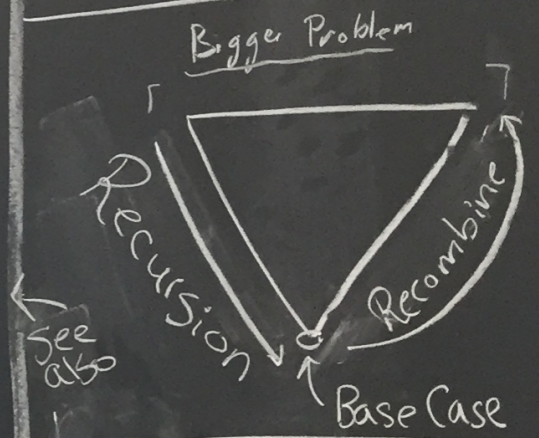
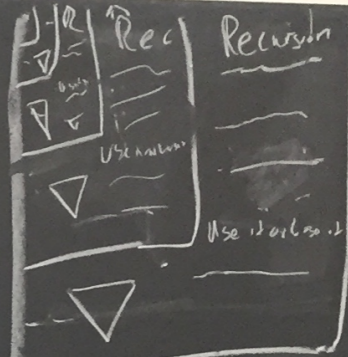
$T(1) = 1 + T(0)$

$T(2) = 1 + T(1)$

$T(N) = N$

which is $O(N)$

*used to find to find time complexity of recursive functions



REPPRESSED MEMORIES

- (1) Align
- (2) Make-change
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- (5) Word difference (EditDistance.py)

Benefi

- Intuitive
- Gets Right
- Useful for structure

8

NAMESPACES ARE ONE HONKING

GREAT IDEA

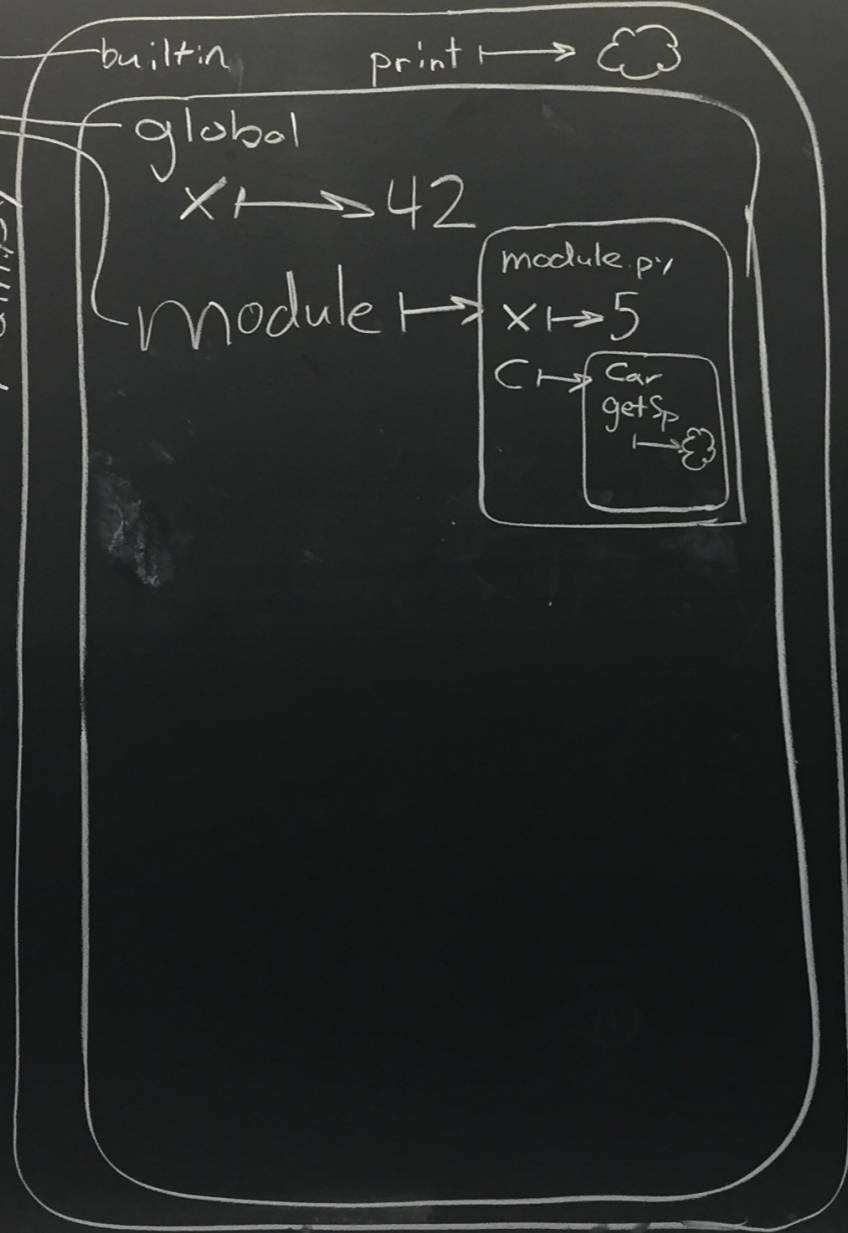
the builtin namespace contains builtin functions

global is for the file

when we import a module, a new namespace gets created

We refer to names in an imported namespace as `module.x`
evaluates to 5

Main.py



How are names looked up?

- 1) Local
- ↳ 2) Global
- ↳ 3) Builtin
- ↳ ERROR

- main.py -
import module

x = 42

```
class Car:
    def __init__(sp):
        self.carSp = sp
    def getSp():
        return self.carSp
```

```
class Honda(Car):
    def __init__(sp):
        super().__init__(sp)
        self.type = "Honda"
    def getType():
        return self.type
```

- module.py -

x = 5

c = Car(10)

h = Honda

s to

s using

to

[n-2]

Dynamic Programming with Tabulation

Why: Recursive functions include a lot of repeat calculations.

What: By keeping track of old calculations you can avoid doing repeat work, making a program run faster.

9

MEMOIZATION:

* When solving a problem, check to see if it has already been computed to avoid redundancy.

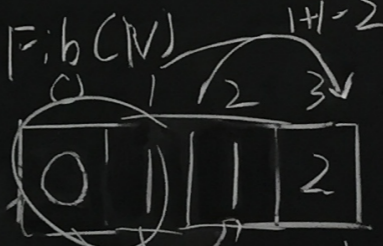
```

fcal if a in memo:
    return memo[a]
else:
    recursion
    
```

Tabulation

* use previous values in a generated table starting at the N=0 case.

Example Fib(N)



Base Case

Recurrence Relation

$$table[n] = table[n-1] + table[n-2]$$

- ① Consider the base case
- ② Calculate how many cells to use
- ③ Use recursion to fill cells using previous cells
- ④ Return result in the result cell

⑩ OOP in a table

Interface

Implementation

- What a thing does
- Determined beforehand

- How it is done (What's in the blackbox)

Public vs Private
(global vs local)

The user uses the public variables and methods (functions).

Java

- Private variables
- Public variables

Python

- "_" notation (private)
- ALL CAPS (global)

* Social Construct!

Objects

- data structures
- classes

Each kind of object has a defined set of operations. The user interacts with instances of objects via their interface.

- new Object (initialization)
- `.this`

- Simply declare the variable ↓
- `self`

- Constructor
- `toString`, `equals`, `hash`

* Knows about itself

• `__init__`

• `__str__`

Inheritance

Subclasses can have more variables and methods available to the user than their superclasses. Each instance of a subclass can work as an instance of its superclass.

- Extends, Implements

- Class Name (superclass)

* Implement the super class

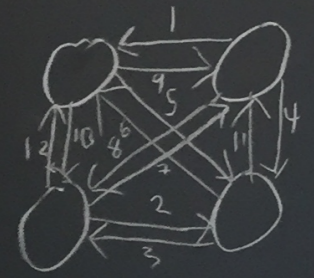
"IS-a"

11

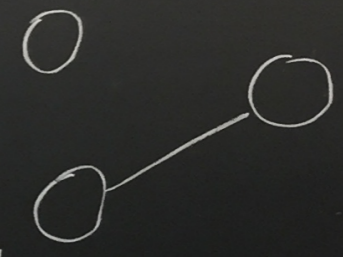
Graphs

Def: a data structure of nodes and edges
- edges connect nodes and can have direction and weight

Uses: modeling things (ie internet)



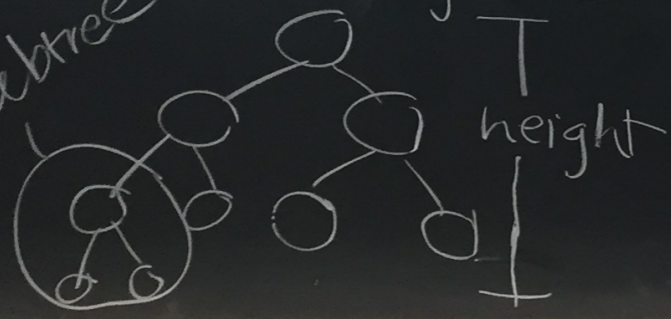
- Types:
- Connected v. nonconnected
 - Complete v. noncomplete
 - directional v. nondirectional
 - weighted v. nonweighted
 - dense v. sparse
 - cyclic v. acyclic



Trees

↳ specific type of connected, directional, nonweighted, acyclic graph
- edges connect roots and leaves

subtree



Binary Search Tree
↳ specific type of tree that makes it easier to store and find things

